

601.220 Intermediate Programming

Summer 2024, Meeting 10 (June 24th)

Today's agenda

- Exercise 15 review
- “Day 17” material
 - Linked lists
 - Exercise 17
- “Day 18” material
 - More linked lists
 - Exercise 18

Reminders/Announcements

- Update on Midterm project deadline: ~~due Friday, June 28th~~
due Saturday, June 29th , 11pm
 - An additional day was added to compensate the fact the starter code was release on Sunday.
 - **No late days allowed** for the midterm project (Please refer to the homework policy at the syllabus tab on the course website)
- Midterm exam: **in class** on Wednesday, July 3th

Exercise 15 review - part 2



- Identify the endianness of ugrad machines
- **Endianness** is the order or sequence in which multi byte words are stored in memory.
 - Little-endian (LE)
 - Big-endian (BE)
- Why is it important to know which system is our computer using?
- Example with base 10 number: 950238851 can be represented in hexadecimal base as 38 A3 7E 83.

Exercise 15 review - part 2

FF \rightsquigarrow 11 11 11 11

Example with base 10 number: 950238851 can be represented in hexadecimal base as 38 A3 7E 83.

sig \leftarrow most byte \rightarrow least sig byte

38	83
A3	7E
7E	A3
83	38

\Downarrow Big endian \Downarrow little endian

Exercise 15 review - part 2



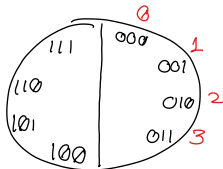
Identify the endianness of ugrad machines

```
(gdb) break endian.c:21
Breakpoint 1 at 0x1243: file endian.c, line 21.
(gdb) run
[...output omitted...]
Breakpoint 1, main () at endian.c:21
21  printf("%u\n", *p);
(gdb) print/x ((unsigned char *)p)[0]
$1 = 0x83
(gdb) print/x ((unsigned char *)p)[1]
$2 = 0x7e
(gdb) print/x ((unsigned char *)p)[2]
$3 = 0xa3
(gdb) print/x ((unsigned char *)p)[3]
$4 = 0x38
```

In base-16, 950238851 is 38A37E83. Since we're seeing the bytes in order from least to most significant, the ugrad machines are *little endian*.

Exercise 15 review - part 3

— — —
↑
negative contribution.



To negate a two's complement value:

- Invert all of the bits (the \sim operator is useful for this)
- Add 1

Exercise 15 review - part 3 $10000\dots$

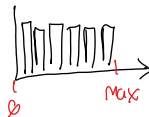
Note that `0x80000000U` is the unsigned int value with only the most significant bit set to 1. This is the sign bit, and values with this bit set are negative.

```
unsigned int magnitude(unsigned int value) {  
    if ((value & 0x80000000U) == 0U) {  
        return value; // value is non-negative  
    }
```

```
    // value is negative, so invert bits and add 1  
    value = ~value; // invert bits  
    value += 1U; // add 1  
    return value;  
}
```

$$\begin{array}{r} 1000\dots \rightarrow 0x80\dots \\ \oplus 1101\dots \rightarrow \text{value} \\ \hline 1000\dots \end{array}$$

Exercise 15 review - part 4



Generating a uniformly distributed pseudo-random integer in the range 0 (inclusive) to max_num (exclusive):

```
int gen_uniform(int max_num) {  
    return rand() % max_num;  
}
```

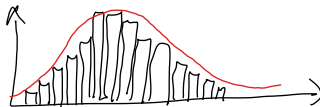
rand() generates integers between 0 and RAND_MAX

Exercise 15 review - part 4

Generating 500 random values in range 0 (inclusive) to `max_range` (exclusive) and tallying them in the `hist` array:

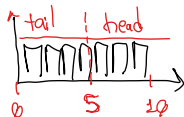
```
for (int i = 0; i < 500; i++) {  
    hist[gen_uniform(max_range)]++;  
}
```

Exercise 15 review



Generating normally-distributed integer values in the range 0 (inclusive) to `max_range` (exclusive):

- Base idea: throw a coin a 1000 times. How many heads would you expect to obtain?
- Now let's repeat this experiment a thousand times. What is the distribution of the total number of heads of all thousand experiments?



Exercise 15 review

Generating normally-distributed integer values in the range 0 (inclusive) to `max_range` (exclusive):

```
int normal_rand(int max_num) {  
    int result = 0;  
    for (int i = 1; i < max_num; i++) {  
        if ((rand() & 1) == 1) {  
            result++;  
        }  
    }  
    return result;  
}
```

This is basically flipping a coin `max_num-1` times and counting how many times it's heads.

Exercise 15 review

Generating 500 normally-distributed values in the range 0 (inclusive) to `max_range` (exclusive) and tallying them in the `hist` array:

```
for (int i = 0; i < 500; i++) {  
    hist[normal_rand(max_range)]++;  
}
```

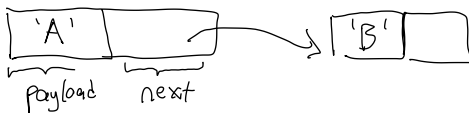
Day 17 recap questions

- 1 Describe the linked list structure by a diagram.
- 2 Compare arrays and linked lists. Write down their pros and cons.
- 3 What is a linked list's head? How is it different from a node? Explain.
- 4 How do you calculate `length` of a linked list?
- 5 How do you implement `add_after` on a singly linked list?

1. Describe the linked list structure by a diagram.

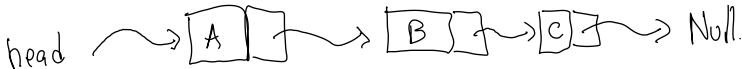
struct Node type:

```
struct Node {  
    char payload; // payload could be any data type  
    struct Node *next;  
};
```



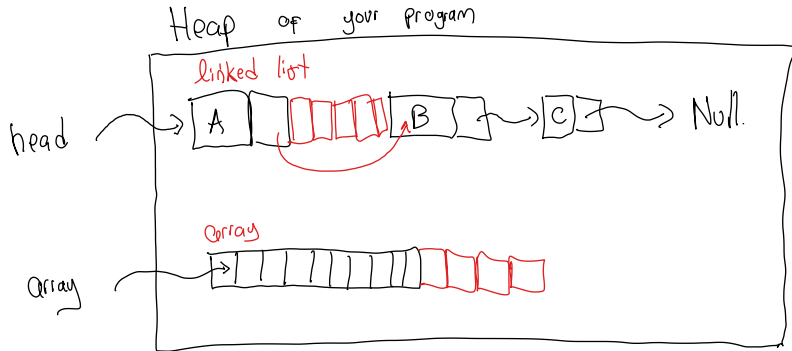
Example linked list

```
// code creating a linked list  
struct Node *head = malloc(sizeof(struct Node));  
head->payload = 'A';  
head->next = malloc(sizeof(struct Node));  
head->next->payload = 'B';  
head->next->next = malloc(sizeof(struct Node));  
head->next->next->payload = 'C';  
head->next->next->next = NULL;
```



A more concise representation

Arrays vs linked lists

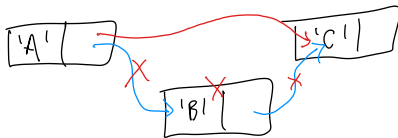


2. Compare arrays and linked lists. Write down their pros and cons.

Arrays:

- Pro: $O(1)$ access to arbitrary element
- Con: $O(N)$ to insert or remove element at arbitrary position
- Pro: better locality (fewer cache misses when iterating)
- Pro: more compact
- Con: fixed size, to reallocate must allocate new array and copy existing data

Linked list pros and cons



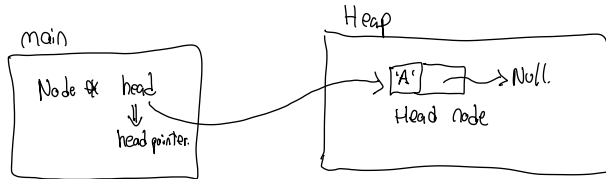
Linked list:

- Con: $O(N)$ access to arbitrary element
- Pro: $O(1)$ to remove element at arbitrary position
- Con: worse locality (more cache misses when iterating)
- Con: less compact (next pointers require space)
- Pro: can grow incrementally, nodes are allocated one at a time

3. What is a linked list's head? How is it different from a node? Explain.

Contrast: *head pointer* vs. *head node*. The head pointer is a pointer variable storing a pointer to the first node. The head node *is* the first node in the linked list.

Picture:



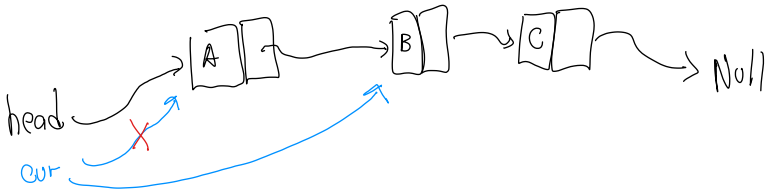
4. How do you calculate length of a linked list?

A loop is required:

```
struct Node *head = /* points to first node */;  
int count = 0;
```

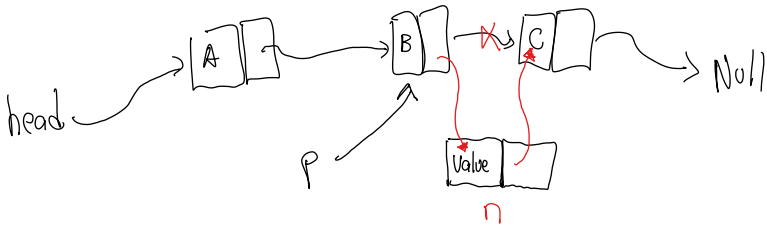
```
for (struct Node *cur = head; cur != NULL; cur = cur->next) {  
    count++;  
}
```

cur → (*cur).next



5. How do you implement add_after on a singly linked list?

```
void add_after(struct Node *p, char value) {  
    struct Node *n = malloc(sizeof(struct Node));  
    n->payload = value;  
    n->next = p->next;  
    p->next = n;  
}
```



Exercise 17

- Basic linked list functions
- Drawing pictures to reason about how linked lists operations should work is very helpful!
- Note that `reverse_print` is most easily implemented using recursion
- Breakout rooms 1–10 are “social”
- Use Slack to let us know if you have questions

Day 18 recap questions

- ① How do you implement *add_front* on a linked list?
- ② How do you modify a singly linked list to create a doubly linked list?
- ③ How do you make a copy of a singly linked list?
- ④ Why does *add_after* takes a struct Node * as input, but *add_front* takes struct Node **?
- ⑤ What cases should be handled when implementing *remove_front*?

Changing the address of a pointer inside a function

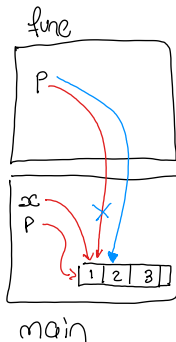
The wrong way of doing it c is a passed by value language!

```
void func(int *p)
{
    p+=1;
    printf("Inside func %d\n", *p);
}
int main()
{
    int x[5] = {1,2,3,4,5};
    int *p = x;
    func(p);
    printf("outside func %d\n", *p);

    return 0;
}
```

Diagram illustrating the state of the program:

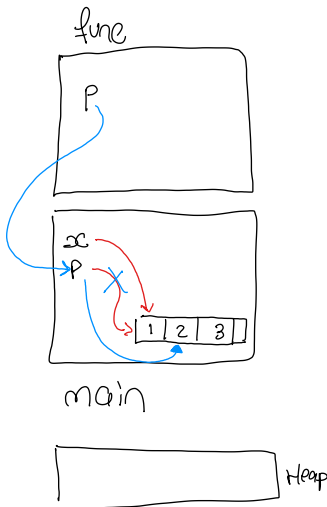
- Red arrow 1 points to the `func` function call in `main`.
- Red arrow 2 points to the `p` parameter in `func`.
- Red arrow 1 points to the `x` array in `main`.



Changing the address of a pointer inside a function

The the right way of doing it

```
void func(int **p)
{
    *p+=1;
    printf("Inside func %d\n", **p);
}
int main()
{
    int x[5] = {1,2,3,4,5};
    int *p = x;
    func(&p);
    printf("outside func %d\n", *p);
}
```



4. Why does `add_after` takes a `struct Node *` as input, but `add_front` takes `struct Node **`?

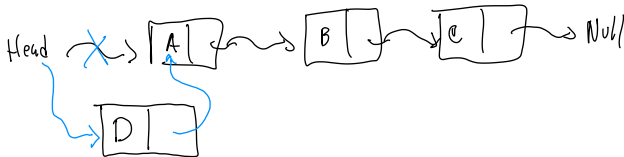
Because `add_after` needs to change which node the head pointer points to. For example:

```
struct Node *head = /* linked list containing 'A', 'B', 'C' */;  
// ...  
add_front(&head, 'D');
```

Before:

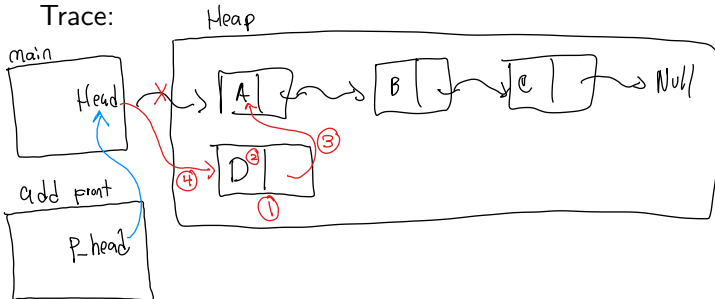


After:



1. How do you implement *add_front* on a linked list?

```
void add_front(struct Node **p_head, char value) {
    struct Node *node = malloc(sizeof(struct Node)); ①
    node->data = value; ②
    node->next = *p_head; ③
    *p_head = node; ④
}
```

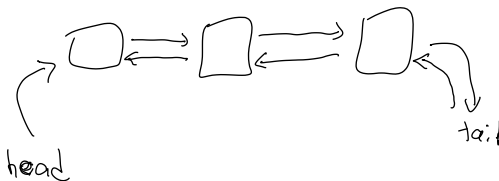


2. How do you modify a singly linked list to create a doubly linked list?

Have each node store a pointer to the *previous* node in the list, in addition to the next node in the list. I.e.:

```
struct Node {  
    char payload;  
    struct Node *prev, *next;  
};
```

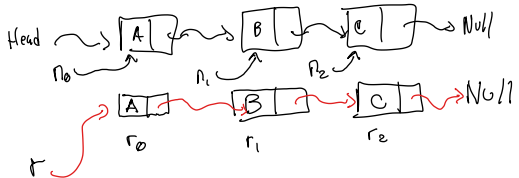
Example:



3. How do you make a copy of a singly linked list?

One way is to use recursion:

```
struct Node *copy_list(struct Node *n) {  
    struct Node *result;  
    if (n == NULL) {  
        result = NULL;  
    } else {  
        result = malloc(sizeof(struct Node));  
        result->payload = n->payload;  
        result->next = copy_list(n->next);  
    }  
    return result;  
}
```



5. What cases should be handled when implementing *remove_front*?

There should not be any special cases.

```
void remove_front(struct Node **p_list) {  
    assert(*p_list != NULL);  
    struct Node *succ = (*p_list)->next;  
    free(*p_list);    // free original head node  
    *p_list = succ;   // make head pointer point to second node  
}
```

Exercise 18

- More linked list operations (including ones requiring pointer to head pointer)
- Again, drawing diagrams is very helpful for reasoning about linked list operations
- Breakout rooms 1–10 are “social”
- Use Slack to let us know if you have any questions!

Notes

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